

# □ LECTURE 5

- PDEs discretization limitations (timestep)
- ODEs
- Physical basis for Moore's scaling & the reason the technical civilization took off like a rocket in the 2<sup>nd</sup> half of 20th century
- Break-down of Moore's law and its consequences for the world of programming
  - **Supercomputing today**
  - Programming and parallelism
  - ODEs

**Literature:** see links on our course home page, the references page, and the coding page available from the course page.

# PDEs

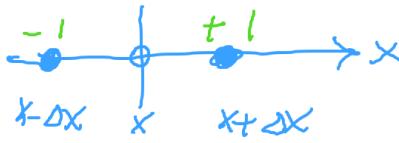
- PDEs use stencils
- Depending on the stencil, there may be limitations on discretization steps
- The implicit schemes often avoid them
- Explicit schemes have timestep and/or spatial discretization limits
- Courant-Friedrichs-Levy is an example for wave-type and hydrodynamics equations

# Von Neumann stability analysis of PDE algorithms: diffusion equation

PDES Example  $u$ -spreading quantity  $u(x, y, z, t)$  (in 3D)  
 $\frac{\partial u}{\partial t} = D \left( \frac{\partial^2 u}{\partial x^2} + \frac{\partial^2 u}{\partial y^2} + \frac{\partial^2 u}{\partial z^2} \right) = D \nabla^2 u = D \Delta u$  diff-coff.:  $D, v$  (alt. names)  
 $[D] = \frac{m^2}{s}$

Spatial differentiation

ex.  $\frac{\partial u}{\partial x} \approx \frac{u(x+\Delta x) - u(x-\Delta x)}{2\Delta x}$



$$\frac{\partial^2 u}{\partial x^2} \approx \frac{1}{(\Delta x)^2} \left[ \frac{u(x+\Delta x) - u(x)}{\Delta x} - \frac{u(x) - u(x-\Delta x)}{\Delta x} \right]$$

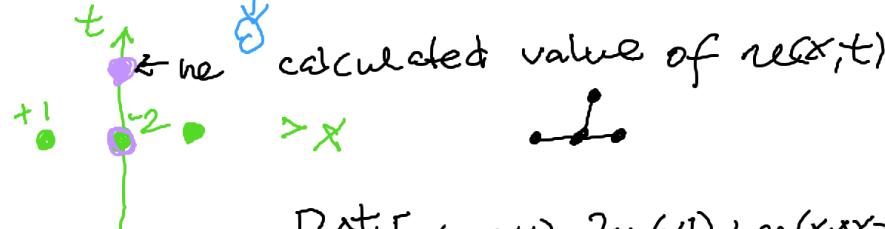
$$= \frac{1}{(\Delta x)^2} (u(x+\Delta x) - 2u(x) + u(x-\Delta x))$$



stencil for heat eq.

(1-D)

$$\frac{\partial u}{\partial t} \approx \frac{u(t+\Delta t) - u(t)}{\Delta t}$$



$$u(x, \dots, t+\Delta t) = \frac{D \Delta t}{(\Delta x)^2} [u(x+\Delta x, t) - 2u(x, t) + u(x-\Delta x, t)] + u(x, t)$$

$$u(x, t+\Delta t) = u(x) + r (u(x) - 2u(x) + u(x-\Delta x))$$

where  $r := \frac{D \Delta t}{(\Delta x)^2}$

# Von Neumann stability analysis of PDE algorithms:

diffusion equation (done below),  
wave equation (exercise for you)

→ Von Neumann stability anal. for PDEs

$$u = u^{\text{num.}} + \text{error } \epsilon$$

$$\text{For linear eqs. } \epsilon: \frac{\partial \epsilon}{\partial t} = D \nabla^2 \epsilon \quad (\text{idea: Fourier})$$

discretize  $u_j^n$  timestep  
 $j \leftarrow$  spatial location

$$\epsilon_j^n:$$

$$\epsilon_j^{n+1} = \epsilon_j^n + r(\epsilon_{j+1}^n - 2\epsilon_j^n + \epsilon_{j-1}^n)$$

$$\epsilon_j^n \sim E(t) e^{ikx}$$

$$e^{ikx}$$

$$\frac{\partial^2 u}{\partial t^2} = c^2 \nabla^2 u$$

wave equation

$\square$  new  $Q: \Delta t \leq ?$

$\Delta t \leq \boxed{\Delta x} \frac{\Delta x}{c}$

Plug that into PDE

$$E(t+\Delta t) e^{ikx} = E(t) e^{ikx} + r(E(t) e^{ik(x+\Delta x)} - E(t) e^{ik(x-\Delta x)}) \leq E(t)$$

Stability condition

$$|E(t+\Delta t)| \leq |E(t)| \Leftrightarrow |1 + r(e^{ikx} + e^{-ikx} - 2)| < 1$$

$$\text{Trig: } \sin \frac{\theta}{2} = \frac{e^{i\theta/2} - e^{-i\theta/2}}{2i} \Rightarrow \sin^2 \frac{\theta}{2} = \left(-\frac{1}{4}\right)(e^{i\theta} + e^{-i\theta} - 2)$$

$$\theta \equiv k\Delta x$$

$$\text{Stab. crit.: } \left| 1 - 4r n \sin^2 \frac{k\Delta x}{2} \right| < 1 \Rightarrow r = \frac{\Delta t}{(\Delta x)^2 n} < \frac{1}{2n}$$

$$n = \# \text{dim}$$

$$\Delta t < \frac{1}{2} \cdot \frac{(kn)^2}{D} \quad (1 \rightarrow)$$

# Special theory of perturbations (numerical calculation)

Most popular numerical integration methods for **ODEs**  
Euler method (1st order) & Runge-Kutta (2nd - 8th order)

Leonard Euler



Carle Runge



1856-1927

Martin Kutta



1867-1944

## The Euler method

We want to approximate the solution of the differential equation

$$y'(t) = f(t, y(t)), \quad y(t_0) = y_0, \quad (1)$$

For instance, the Kepler problem which is a 2nd-order equation, can be turned into the 1st order equations by introducing double the number of equations and variables: e.g., instead of handling the second derivative of variable  $x$ , as in the Newton's equations of motion, one can integrate the first-order (=first derivative only) equations using variables  $x$  and  $v_x = dx/dt$  (that latter definition becomes an additional equation to be integrated).

Starting with the differential equation (1), we replace the derivative  $y'$  by the finite difference approximation, which yields the following formula

which yields 
$$y'(t) \approx \frac{y(t+h) - y(t)}{h}, \quad (2)$$

$$y(t+h) \approx y(t) + h f(t, y(t)) \quad (3).$$

This formula is usually applied in the following way.

## The Euler method (cont'd)

$$y(t + h) \approx y(t) + hf(t, y(t)) \quad (3).$$

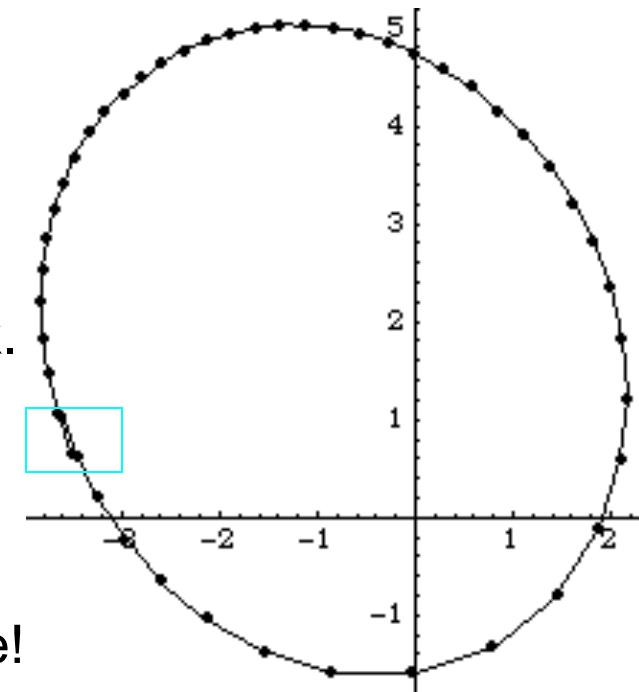
This formula is usually applied in the following way.

We choose a step size  $h$ , and we construct the sequence  $t_0$ ,  $t_1 = t_0 + h$ ,  $t_2 = t_0 + 2h$ , ... We denote by  $y_n$  a numerical estimate of the exact solution  $y(t_n)$ . Motivated by (3), we compute these estimates by the following recursive scheme

$$y_{n+1} = y_n + h f(t_n, y_n).$$

This is the *Euler method* (1768), probably invented but not formalized earlier by Robert Hook.

It's a first (or second) order method, meaning that the total error is  $\sim h^1$ <sup>(2)</sup> It requires small time steps & has only moderate accuracy, but it's very simple!



## The classical fourth-order Runge-Kutta method

One member of the family of Runge-Kutta methods is so commonly used, that it is often referred to as "RK4" or simply as "*the* Runge-Kutta method".

The RK4 method for the problem

$$y'(t) = f(t, y(t)), \quad y(t_0) = y_0, \quad (1)$$

is given by the following equation:

$$y_{n+1} = y_n + \frac{h}{6}(k_1 + 2k_2 + 2k_3 + k_4)$$

where

$$k_1 = f(t_n, y_n)$$

$$k_2 = f\left(t_n + \frac{h}{2}, y_n + \frac{h}{2}k_1\right)$$

$$k_3 = f\left(t_n + \frac{h}{2}, y_n + \frac{h}{2}k_2\right)$$

$$k_4 = f(t_n + h, y_n + hk_3)$$

Thus, the next value ( $y_{n+1}$ ) is determined by the present value ( $y_n$ ) plus the product of the size of the interval ( $h$ ) and an estimated slope.

## Runge-Kutta 4th order (cont' d)

Thus, the next value ( $y_{n+1}$ ) is determined by the present value ( $y_n$ ) plus the product of the size of the interval ( $h$ ) and an estimated slope

$$\text{slope} = \frac{k_1 + 2k_2 + 2k_3 + k_4}{6}.$$

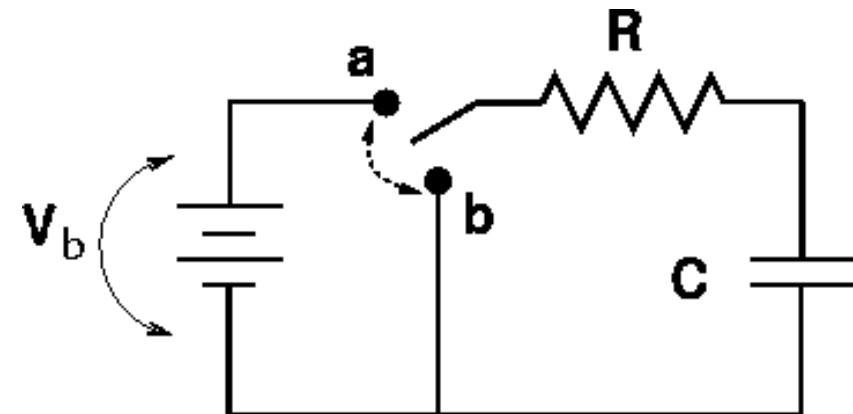
Thus, the next value ( $y_{n+1}$ ) is determined by the present value ( $y_n$ ) plus the product of the size of the interval ( $h$ ) and an estimated slope. The slope is a weighted average of slopes:

- $k_1$  is the slope at the beginning of the interval;
- $k_2$  is the slope at the midpoint of the interval, using slope  $k_1$  to determine the value of  $y$  at the point  $t_n + h/2$  using Euler's method;
- $k_3$  is again the slope at the midpoint, but now using the slope  $k_2$  to determine the  $y$ -value;
- $k_4$  is the slope at the end of the interval, with its  $y$ -value determined using  $k_3$ .

When the four slopes are averaged, more weight is given to the midpoint.

The RK4 method is a fourth-order method, meaning that the total error is  $\sim h^4$ . It allows larger time steps & better accuracy than low-order methods.

# Rise and present slow-down of the exponential growth of our TECHNICAL CIVILIZATION

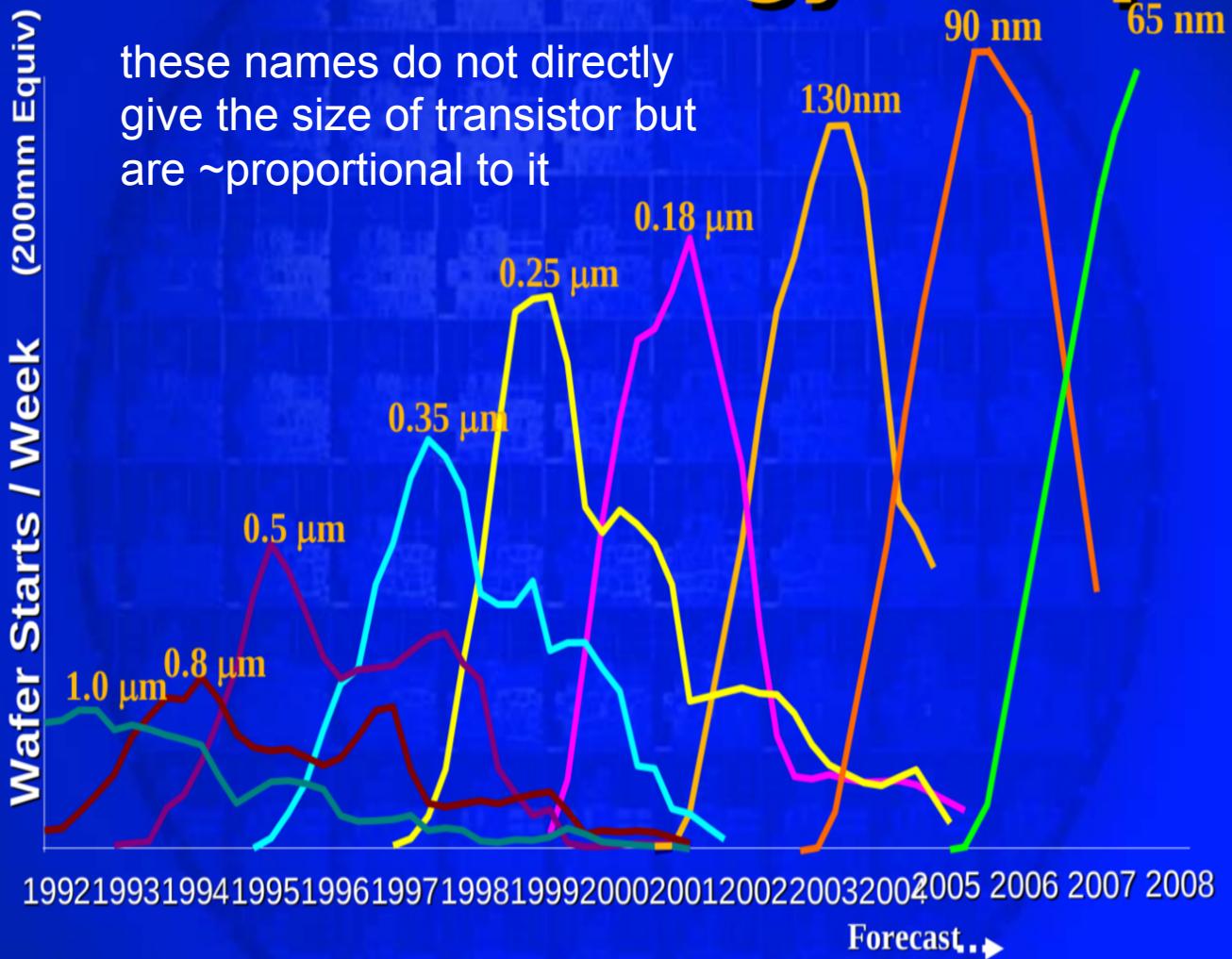


~10,000,000,000 transistors, acting like RC circuits, are etched in every CPU and GPU chip of a modern workstation ( $10^7$  -  $10^8$  /mm $^2$ )



# Intel Technology Ramps

these names do not directly give the size of transistor but are ~proportional to it

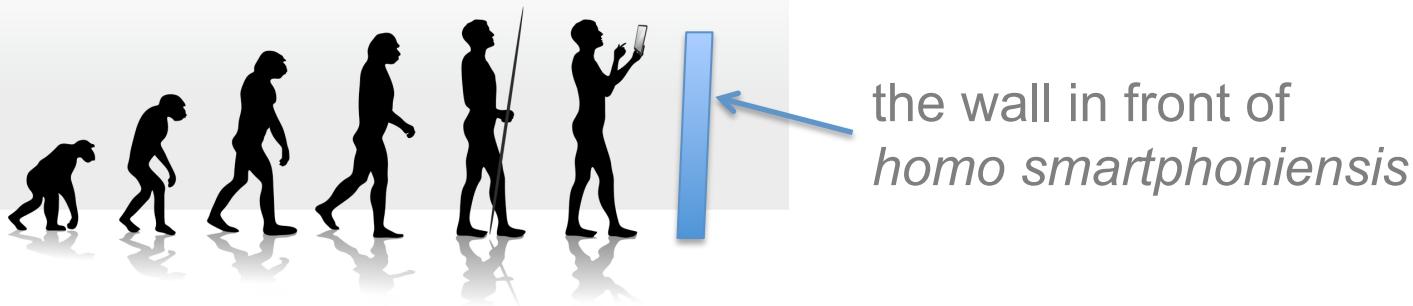


New Technology Ramps Every 2 Years

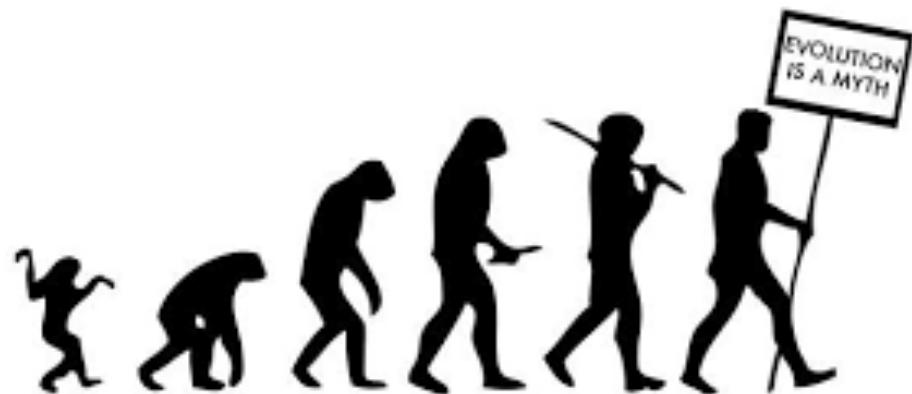
Name of technology cycle:

- 10  $\mu\text{m}$  – 1971
- 6  $\mu\text{m}$  – 1974
- 3  $\mu\text{m}$  – 1977
- 1.5  $\mu\text{m}$  – 1981
- 1  $\mu\text{m}$  – 1984
- 800 nm – 1987
- 600 nm – 1990
- 350 nm – 1993
- 250 nm – 1996
- 180 nm – 1999
- 130 nm – 2001
- 90 nm – 2003
- 65 nm – 2005
- 45 nm – 2007
- 32 nm – 2009
- 22 nm – 2012
- 14 nm – 2014
- 10 nm – 2016
- 7 nm – 2019
- 5 nm – 2020+

Current density of transistors is of order 100 MT/mm<sup>2</sup>

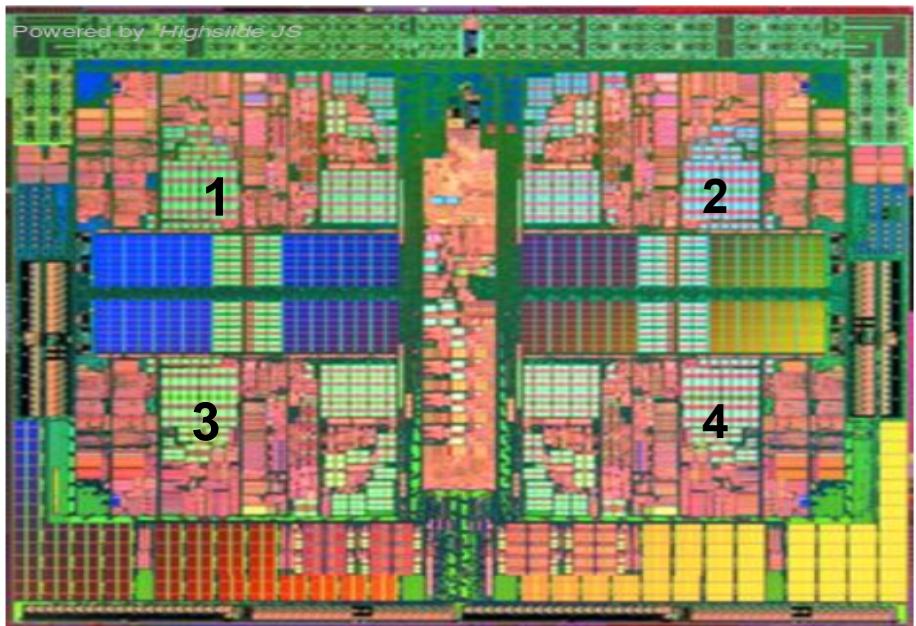
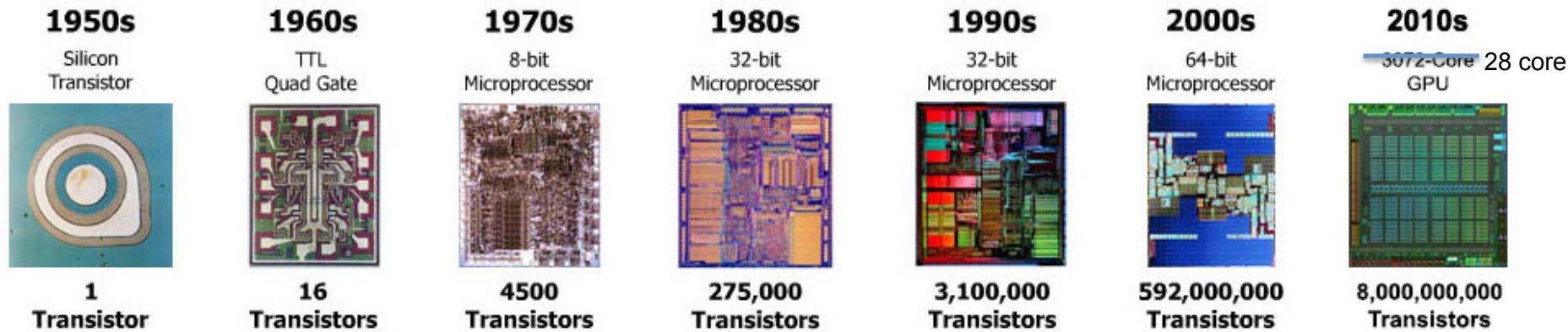


## RISE & CURRENT SLOWDOWN of the exponential growth of our TECHNICAL CIVILIZATION



# 1970s—2010s, the integrated circuit revolution

as a basis of **all our current technical civilization**: 40+ years of **Moore's law**

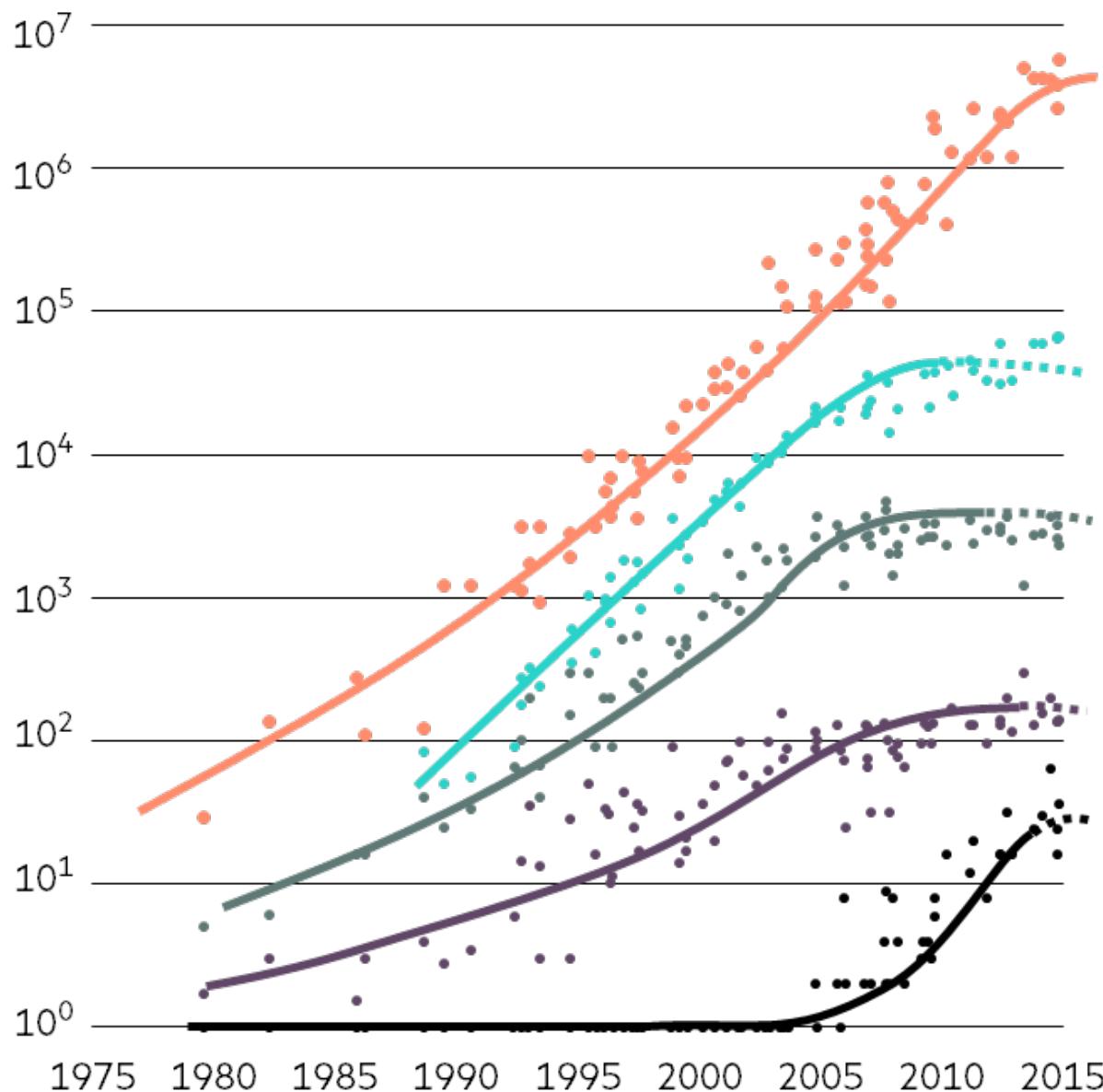


A highly integrated circuit (IC) has both arithmetic + logic units (ALU), fast cache memory hierarchy, and communication circuits, all in a small package of a Central Processing Unit (CPU).

$$N = \text{up to } 10^{10} \text{ transistors}$$

This processor looks like a 4-core CPU or old GPU (Central or Graphics Processing Unit). It is not easy to execute one calculation simultaneously on many **cores** (~ sub-CPUs) of a CPU. We call such processors **multicore**, and programs running on all cores **multithreaded**.

# Microprocessors



after ~2004:

Transistors (thousands) still goes up, a bit slower

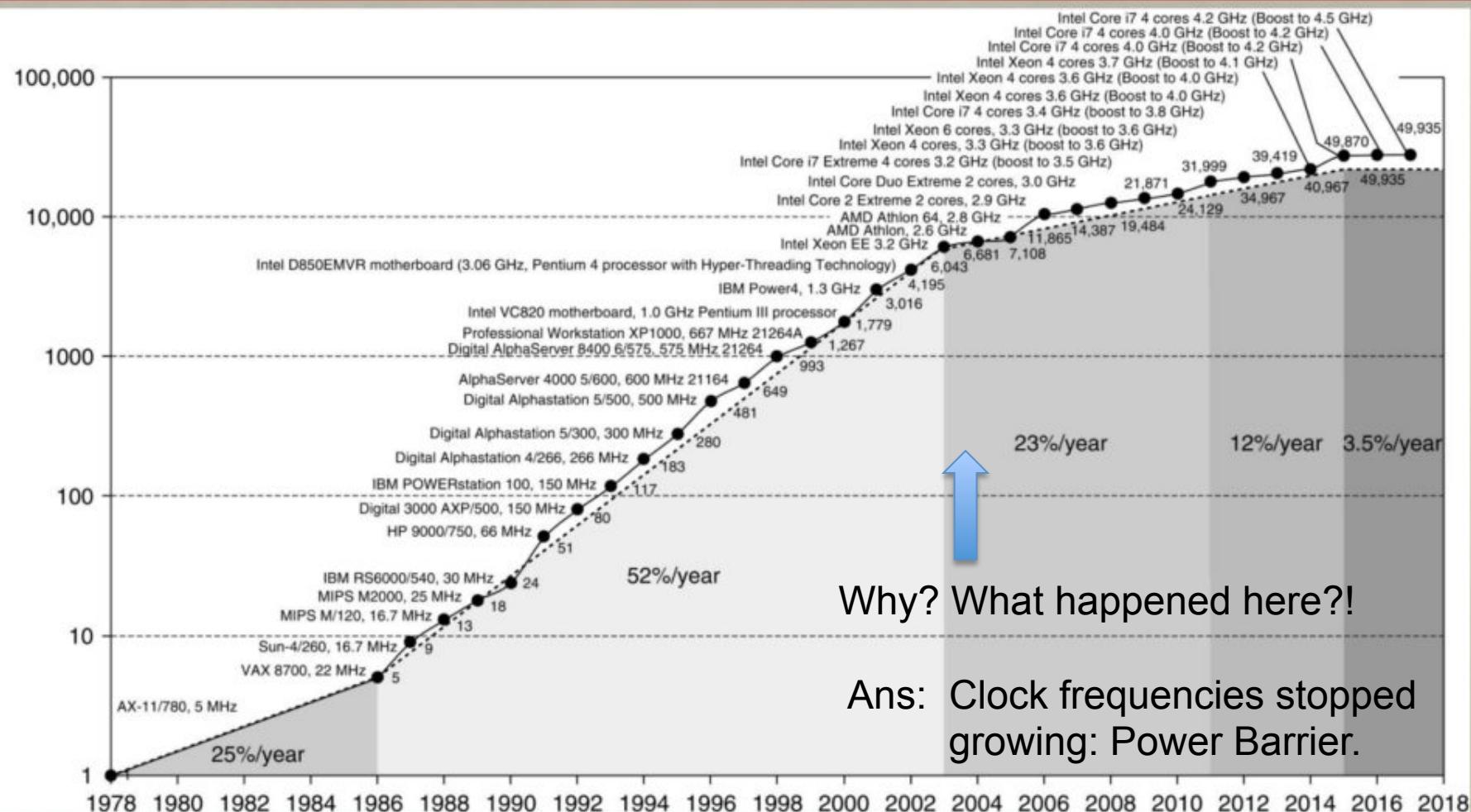
Single-thread Performance (SpecINT) stagnated

Frequency (MHz) stagnated

Typical Power (Watts) hit a wall

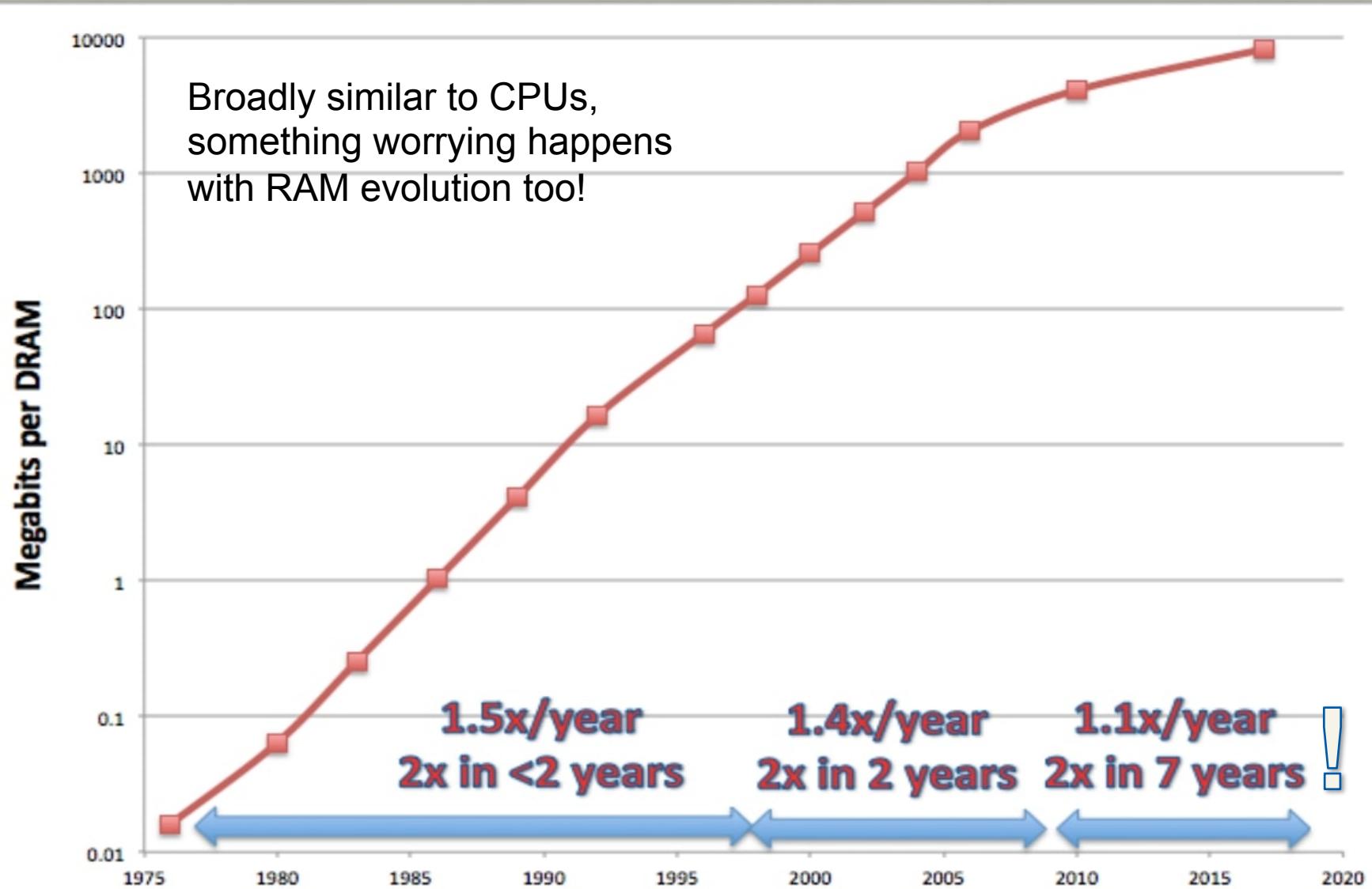
Number of Cores started rapidly climbing

# UNIPROCESSOR PERFORMANCE (SINGLE CORE)



# MOORE'S LAW IN DRAMs (memory)

<https://www.nextbigfuture.com/2019/02/the-end-of-moores-law-in-detail-and-starting-a-new-golden-age.html> = a DARPA talk



**Top 500 supercomputers in the world  
since ~2008 grow in speed slower than before**

